# The van der Corput Embedding of *ax* + *b* and its Interval Exchange Map Approximation

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### 1. Introduction

In [4], we introduce interval preserving map approximation of a linear map 3x+1, to attack the well-known and still unsolved 3x + 1 problem, which is firstly proposed by Lothar Collatz in 1930's:

**Conjecture 1.1.** Consider a map  $f : \mathbf{N} \to \mathbf{N}$  such that

$$f(n) = \begin{cases} 3n+1, & \text{if } n \text{ is odd,} \\ n/2, & \text{if } n \text{ is even.} \end{cases}$$

Then for each natural number n, there exists a finite number t such that  $f^t(n) = \underbrace{f \circ f \circ \cdots \circ f(n)}_{t-times} = 1.$ 

In the investigations [3][4], we use the binary van der Corput embedding  $\beta$  of natural numbers into unit interval [0, 1], which gives a low-discrepancy sequence over [0, 1], firstly introduced by van der Corput in 1935(cf.[7] pp 129), and construct a right continuous bijection F over [0, 1) conjugate to 3x + 1 by way of  $\beta$ . In this note, we investigate a right continuous bijection F over [0, 1) conjugate to the linear function ax + b, its 'finite bit approximations' which bring dynamics of interval exchange maps on [0, 1), and observe substitution dynamics and transducers associated with the approximations. As a result, we obtain a necessary and sufficient condition for the minimality of the dynamics ([0, 1), F) (Theorem 3.8).

#### 2. The binary van der Corput embedding of natural numbers and a conjugacy of ax + b

We follow the notations in [4].

**Definition 2.1.** Let  $n = g_k \cdot 2^k + g_{k-1} \cdot 2^{k-1} + \cdots + g_0$  (=  $(g_k g_{k-1} \cdots g_0)_2$  for short) be a binary expansion of a natural number n. Then we define  $\beta : \mathbf{N} \to [0, 1)$ , given by

$$\beta(n) = \frac{g_0}{2} + \frac{g_1}{2^2} + \dots + \frac{g_k}{2^{k+1}} = (0.g_0g_1\cdots g_k)_2$$

As  $\beta : \mathbf{N} \to [0,1)$  is one-to-one and  $\beta$  embeds natural numbers densely into [0,1), we call  $\beta$  the binary van der Corput embedding.

For a binary expression  $n = (g_l g_{l-1} \cdots g_0)_2$ , we put  $l = \operatorname{ord}(n) := [\log_2 n]$  where [x] stands for an integer not greater than x, and  $n|^k$  denotes an *upper cut off* of n at k-th order;

$$n|^k = (g_{k-1}g_{k-2}\cdots g_0)_2 \equiv n \pmod{2^k}.$$

To make free from ambiguity, we always take the finite binary expression  $\beta(n) = (0.g_0g_1\cdots g_l)_2$  of any natural number n. For a real number  $x \in [0, 1)$  and  $k \in \mathbb{N}$ ,  $x|_k$  denotes a cut off of x at (-k-1)-th order in the binary expression;

$$x|_k = (0.e_1e_2\cdots e_k)_2 = \sum_{j=1}^k \frac{e_j}{2^j}, \text{ if } x = (0.e_1e_2\cdots)_2 = \sum_{j=1}^\infty \frac{e_j}{2^j}.$$

We use the notation  $[x]_k$  as a segment  $[x|_k, x|_k + \frac{1}{2^k})$ , hence a natural decomposition of segments

(2.1) 
$$[ (0.e_1e_2\cdots e_k)_2 )_k = [ (0.e_1e_2\cdots e_k0)_2 )_{k+1} \oplus [ (0.e_1e_2\cdots e_k1)_2 )_{k+1} ,$$

where  $[a, b] \oplus [b, c)$  stands for a division of an interval [a, c) at b. Then, the k-th refinement of the unit interval [0, 1) is given as

(2.2) 
$$[0,1) = \bigoplus_{j=0}^{2^{k}-1} \left[ \beta(\nu_{j}^{(k)}) \right]_{k},$$

where  $\{\nu_j^{(k)} \mid j = 0, \dots, 2^k - 1\} = \{0, 1, \dots, 2^k - 1\}$  with  $\beta(\nu_i^{(k)}) < \beta(\nu_j^{(k)})$  if and only if i < j. By definition, we easily to see the followings.

**Lemma 2.2.** For  $x, y \in [0, 1)$  and  $k, l, m, n \in \mathbf{N}$ ,

- (1)  $x \in [x]_k$ . (2)  $[x]_k \cap [y]_k \neq \emptyset$  holds if and only if  $x|_k = y|_k$ , and hence  $[x]_k = [y]_k$ . (3) if  $x \neq y$ , there exists  $l \in \mathbf{N}$  such that  $[x]_l \cap [y]_l = \emptyset$ . (4)  $[x]_{k+1} \subsetneq [x]_k$  and  $\bigcap_{k=1}^{\infty} [x]_k = \{x\}.$ (5)  $\beta(m)|_k = \beta(m|^k)$ , thus  $[\beta(m)]_k = [\beta(n)]_k$  if and only if  $m|^k = n|^k$ . (6) If l > ord(m), then  $\beta(m+n \cdot 2^l) = \beta(m) + \frac{\beta(n)}{2^l}$ .
- (7)  $(n|^k)|^l = n|^{\min\{k,l\}}.$

Let us consider a linear function f(n) = an + b on natural numbers, where  $a, b \in \mathbf{N}$  and a > b. We are to construct the conjugacy  $F: [0,1) \to [0,1)$  of  $f: \mathbf{N} \to \mathbf{N}$ , that is,  $F \circ \beta(n) = \beta \circ f(n)$  holds for any natural number n.

Lemma 2.3. Suppose a is odd number. Then we have

- (1)  $[\beta \circ f(m)]_k = [\beta \circ f(n)]_k$  holds if and only if  $[\beta(m)]_k = [\beta(n)]_k$  for any  $m, n \in \mathbb{N}$ .
- (2)  $\{\beta \circ f(n) \mid n \in \mathbf{N}\}$  is dense in [0, 1).

*Proof.* (1) By Lemma 2.2 (5),  $[\beta(m)]_k = [\beta(n)]_k$  means  $m \equiv n \pmod{2^k}$ , which equivalents to  $am + b \equiv b$  $an + b \pmod{2^k}$  as a and  $2^k$  are coprime, hence  $[\beta(am + b)]_k = [\beta(an + b)]_k$ .

(2) Given  $x \in [0,1)$  and  $k \in \mathbf{N}$ , consider  $m_k = \beta^{-1}(x|_k)$ , and take  $b' \in \mathbf{N}$  with  $b' \equiv b - m_k \pmod{a}$ . As a and  $2^k$  is coprime, we have  $2^{k\varphi(a)} = (2^k)^{\varphi(a)} \equiv 1 \pmod{a}$ , hence  $m_k + b' \cdot 2^{k\varphi(a)} \equiv m_k + b' \equiv b \pmod{a}$ , where  $\varphi$  stands for Euler's totient function. Thus there exists  $n \in \mathbf{N}$  with  $m_k + b' \cdot 2^{k\varphi(a)} = an + b$ , and hence

$$\beta(an+b) = \beta(m_k + b' \cdot 2^{k\varphi(a)}) = x|_k + \frac{\beta(b')}{2^{k\varphi(a)}} \in [x]_k.$$
  
[0,1) and  $k \in \mathbf{N}$ , there exists  $n \in \mathbf{N}$  with  $|x - \beta(an+b)| < 2^{-k}$ .

As a result, for any  $x \in [0,1)$  and  $k \in \mathbf{N}$ , there exists  $n \in \mathbf{N}$  with  $|x - \beta(an + b)| < 2^{-k}$ .

Now we suppose a is odd. Then the conjugacy F is defined on the dense set  $\beta(\mathbf{N})$  by  $F(\beta(n)) := \beta(f(n))$ , and its extension to [0,1) is given as follows. For any  $x \in [0,1)$ , consider a sequence  $n_k = \beta^{-1}(x|_k)$ ,  $k = 1, 2, \ldots$  Obviously we see  $n_{k+1} \equiv n_k \pmod{2^k}$ , hence  $an_{k+1} + b \equiv an_k + b \pmod{2^k}$ . It follows from Lemma 2.2 (4) that

$$\left[\beta(an_{k+1}+b)\right]_{k+1} \subsetneq \left[\beta(an_{k+1}+b)\right]_{k} = \left[\beta(an_{k}+b)\right]_{k},$$

holds, and then  $\bigcap_{k \to \infty}^{\infty} [\beta(an_k + b))_k$  consists of a unique point, denoted by  $\lim_{k \to \infty} \beta(an_k + b)$ . Thus we define the conjugacy  $F: [0,1) \to [0,1)$  of the linear function f:

(2.3) 
$$F(x) = \lim_{k \to \infty} \beta \left( a\beta^{-1}(x|_k) + b \right)$$

**Lemma 2.4.** For  $x, y \in [0, 1)$ , we have

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- (1)  $[F(x|_k)]_k = [F(x)]_k$ , hence  $F(x)|_k = F(x|_k)|_k$ .
- (2)  $F([x]_k) \subset [F(x)]_k$ , hence F is right continuous.
- (3)  $[F(x)]_k = [F(y)]_k$  if and only if  $[x]_k = [y]_k$ , hence F is injective.

*Proof.* (1) By definition, we have

$$\{F(x)\} = \bigcap_{k=1}^{\infty} \left[ \beta(a\beta^{-1}(x|_k) + b) \right]_k \subset \left[ \beta(a\beta^{-1}(x|_k) + b) \right]_k = \left[ F(x|_k) \right]_k,$$

that is,  $F(x) \in [F(x|_k))_k$ , while  $F(x) \in [F(x))_k$  by Lemma 2.2 (1). Thus we see  $[F(x|_k))_k \cap [F(x))_k \neq \emptyset$ , which means  $[F(x|_k))_k = [F(x))_k$  by Lemma 2.2 (2).

(2)  $y \in [x]_k$  means  $y|_k = x|_k$ . It follows from (1) that

$$F(y) \in \left[ \begin{array}{c} F(y) \end{array} \right]_k = \left[ \begin{array}{c} F(y|_k) \end{array} \right]_k = \left[ \begin{array}{c} F(x|_k) \end{array} \right]_k = \left[ \begin{array}{c} F(x) \end{array} \right]_k$$

For any  $\varepsilon > 0$ , take  $k \in \mathbf{N}$  with  $2^{-k} < \varepsilon$ . Then for any y with  $x \le y \le x + 2^{-(k+1)}$ , we see  $y \in [x]_k$  and

$$F(y) \in F([x]_k) \subset [F(x)]_k,$$

that is,  $|F(y) - F(x)| < 2^{-k} < \varepsilon$ , showing the right continuity of F.

(3) Put  $m_k = \beta^{-1}(x|_k)$  and  $n_k = \beta^{-1}(y|_k)$ , then we see  $[F(x)]_k = [F(x|_k)]_k = [\beta \circ f(m_k)]_k$  and  $[F(y)]_k = [F(y|_k)]_k = [\beta \circ f(n_k)]_k$  by (1). It follows from Lemma 2.3 (1) that  $[F(x)]_k = [F(y)]_k$  if and only if  $[x]_k = [\beta(m_k)]_k = [\beta(n_k)]_k = [y]_k$ . Suppose that  $x \neq y$ , then  $[x]_k \cap [y]_k = \emptyset$  holds for some k. Then we have  $[F(x)]_k \cap [F(y)]_k = \emptyset$ , hence  $F(x) \neq F(y)$ .

Lemma 2.4 (1) is an extension of a natural property

(2.4) 
$$f(n|^k)|^k = f(n)|^k$$

of the linear function f(n) = an + b with  $a, b \in \mathbf{N}$ , while (2)(3) implies that F brings segment-wise exchange over the refinement (2.2) of the unit interval in any order k. Indeed we have the following.

**Proposition 2.5.** For any  $k \in \mathbf{N}$  and  $x \in [0,1)$ , the conjugacy F gives a right continuous bijection

$$F:[x]_k \to [F(x)]_k.$$

Particularly, F preserves the Lebesgue measure  $\mu$  on [0,1).

*Proof.* All we have to show is surjectivity of F. For any  $y \in [0,1)$  and  $l \in \mathbf{N}$ , there exists  $n_l \in \mathbf{N}$  with  $F \circ \beta(n_l) = \beta \circ f(n_l) \in [y]_l$  by Lemma 2.3 (2). By Lemma 2.4, we see

$$F([\beta(n_l))_l) \subset [F(\beta(n_l)))_l = [y]_l.$$

Let x be the unique accumulation point  $\{x\} = \bigcap_{l=1}^{\infty} [\beta(n_l)]_l$ . Then  $F(x) \in [F(\beta(n_l))]_l \subset [y]_l$  holds for any l meaning that

l, meaning that

$$F(x)\in \bigcap_l^\infty [\ y\ )_l=\{y\},$$

hence y = F(x) and F is a surjection  $[0,1) \to [0,1)$ . Combining Lemma 2.4 (2),(3) and surjectivity of F, we see  $F([x]_k) = [F(x)]_k$  for any  $k \in \mathbf{N}$ . Note that  $\mu([x]_k) = \mu([F(x)]_k) = 2^{-k}$ . It is shown that the  $\sigma$ -field generated by the segments  $[x]_k, k \in \mathbf{N}, x \in [0,1)$  coincides with the Lebesgue measurable sets. Hence the assertion.

Note that for any natural number n,  $[\beta(n)]_k = [\beta(n), \beta(n) + 2^{-k})$  if  $k > \operatorname{ord}(n)$ , thus  $w \in [\beta(n)]_k$  means  $w \ge \beta(n)$ . Therefore we state the right continuity of F only. Indeed, for f(n) = n + 1, we see

$$\lim_{\substack{w \to (0.11)_2 \\ w < (0.11)_2}} F(w) = (0.1)_2 \neq (0.001)_2 = F((0.11)_2).$$

#### 3. Segment-wise linear approximation of the conjugacy F

In view of Proposition 2.5, we construct a sequence of segment-wise linear functions  $F_k$ 's approximating the conjugacy F.

**Definition 3.1.** For each  $k \in \mathbf{N}$ , we define the k-th approximant  $F_k$  as

$$F_k(x) = x - x|_k + F(x)|_k.$$

If  $x \in [\beta(n)]_k$ , we see

$$F_k(x) = x - \beta(n)\big|_k + F(\beta(n))\big|_k,$$

thus  $F_k([\beta(n))_k) = [F \circ \beta(n))_k$ , compatible with Proposition 2.5. Moreover, the fact  $F(x), F_k(x) \in [F \circ \beta(n))_k$  for any  $x \in [\beta(n))_k$  means

$$|F(x) - F_k(x)| < 2^{-k}$$

for any  $x \in [\beta(n)]_k$ , and then for any  $x \in [0,1)$ . Therefore the sequence  $F_k$ , k = 1, 2, ... approximates F uniformly on [0,1), so  $F_k$  simulates the behavior of F that exchange the segments  $[\beta(n)]_k$ 's. We also note that (2.3) is expressed as  $F(x) = \lim_{k \to \infty} F_k(x)$ .

In the following subsections, we investigate the approximants  $F_k$ 's to extract dynamical characteristics of the map f(n) = an + b.

3.1. Behavior of carries and exchange of segments. For each  $k \in \mathbf{N}$ , we define an integer valued function

$$au_k(n) = \left[\frac{f(n)^k}{2^k}\right] = \left[\frac{an^{k+b}}{2^k}\right],$$

then we see

(3.1) 
$$f(n)^{k} = \tau_{k}(n) \cdot 2^{k} + f(n)^{k}.$$

Namely, the function  $\tau_k$  describes the amount of the carry from the lower k bits in the calculation  $n \mapsto an+b$ .

**Theorem 3.2.** Suppose that a is odd and  $0 \le b < a$ . Then, for  $n, k \in \mathbf{N}$ , we have

(1)  $\tau_k(n) \in \{0, 1, \dots, a-1\}.$ 

(2) For binary expressions  $n = (g_p \cdots g_0)_2$  and  $f(n) = (h_q \cdots h_0)_2$ , we have

$$h_k = (g_k + \tau_k(n)) |^1, \ k = 1, \dots, q$$

and

$$\tau_{k+1}(n) = \left[\frac{ag_k + \tau_k(n)}{2}\right],$$

where  $g_k = 0$  for k > p.

(3) Let  $(0.g_0g_1\cdots g_p)_2$  be the binary expression of  $\beta(n)$ . Then, the image of the segment

$$\beta(n) \rangle_{k} = [ (0.g_{0} \cdots g_{k-1})_{2} \rangle_{k} = [ (0.g_{0} \cdots g_{k-1})_{2} \rangle_{k+1} \oplus [ (0.g_{0} \cdots g_{k-1})_{2} \rangle_{k+1}$$

 $by \ F \ is \ given \ as$ 

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$$F\left(\left[\beta(n)\right]_{k}\right) = \begin{cases} \left[F((0.g_{0}\cdots g_{k-1}0)_{2})\right]_{k+1} \oplus \left[F((0.g_{0}\cdots g_{k-1}1)_{2})\right]_{k+1}, & \text{if } \tau_{k}(n) \text{ is even,} \\ \left[F((0.g_{0}\cdots g_{k-1}1)_{2})\right]_{k+1} \oplus \left[F((0.g_{0}\cdots g_{k-1}0)_{2})\right]_{k+1}, & \text{if } \tau_{k}(n) \text{ is odd.} \end{cases}$$

*Proof.* (1) As  $n|^k \leq 2^k - 1$  and a > b, we have

$$0 \le \left[\frac{an|^k + b}{2^k}\right] \le \left[\frac{a \cdot 2^k + b - a}{2^k}\right] = \left[a - \frac{a - b}{2^k}\right] \le a - 1.$$

(2) By definition, we see  $n|^{k+1} = (g_k g_{k-1} \cdots g_0)_2 = g_k \cdot 2^k + n|^k$ . We also see  $f(n)|^{k+1} = h_k \cdot 2^k + f(n)|^k$ , while, by (3.1), we have

(3.2) 
$$f(n)\Big|^{k+1} = f(n|^{k+1})\Big|^{k+1} = \left(a(g_k \cdot 2^k + n|^k) + b\right)\Big|^{k+1} = \left(ag_k \cdot 2^k + f(n|^k)\right)\Big|^{k+1} \\ = \left(\left(ag_k + \tau_k(n)\right) \cdot 2^k + f(n)\Big|^k\right)\Big|^{k+1} = \left(ag_k + \tau_k(n)\right)\Big|^1 \cdot 2^k + f(n)\Big|^k.$$

As a is odd, we obtain the equation  $h_k = (ag_k + \tau_k(n))|^1 = (g_k + \tau_k(n))|^1$ . Again by (3.1) and  $f(n)|^{k+1} = h_k \cdot 2^k + f(n)|^k$ , we see

$$f(n|^{k+1}) = \tau_{k+1}(n) \cdot 2^{k+1} + f(n)|^{k+1} = (2\tau_{k+1}(n) + h_k)2^k + f(n)|^k$$

while, as  $n|^{k+1} = g_k \cdot 2^k + n|^k$ , we have

$$f(g_k \cdot 2^k + n|^k) = ag_k \cdot 2^k + f(n|^k) = (ag_k + \tau_k(n)) \cdot 2^k + f(n)|^k,$$

hence  $2\tau_{k+1}(n) + h_k = ag_k + \tau_k(n)$ . Since  $0 \le h_k/2 < 1$  and  $\tau_{k+1}(n)$  is an integer, we come to

$$\tau_{k+1}(n) = \left[\tau_{k+1}(n) + \frac{h_k}{2}\right] = \left[\frac{ag_k + \tau_k(n)}{2}\right].$$

(3) It comes from (3.2) that

$$\beta(f(n))\big|_{k+1} = \beta(f(n))\big|_k + \frac{1}{2^{k+1}}\beta(g_k + \tau_k(n))\big|_1.$$

Then we have  $(g_k + \tau_k(n))|^1 = g_k$  whenever  $\tau_k(n)$  is even, and  $(g_k + \tau_k(n))|^1 = 1 - g_k$  whenever  $\tau_k(n)$  is odd. Thus we see

$$\left(F \circ \beta(n|^{k} + 2^{k})\right)\Big|_{k+1} = \begin{cases} \left(F \circ \beta(n)\right)\Big|_{k} + \frac{1}{2^{k+1}}, & \text{if } \tau_{k}(n) \text{ is even,} \\ \left(F \circ \beta(n)\right)\Big|_{k} - \frac{1}{2^{k+1}}, & \text{if } \tau_{k}(n) \text{ is odd,} \end{cases}$$

hence the assertion.

## 3.2. Substitution dynamics induced by the conjugacy F.

**Definition 3.3.** For each approximation order k, we label the segment  $[\beta(n)]_k$  as

$$L_k(n) := \tau_0(n)\tau_1(n)\cdots\tau_k(n).$$

We say the segment  $[\beta(n)]_k$  is painted in color  $\tau_k(n)$ .

By Theorem 3.2 (1), the label  $L_k(n)$  is a string over the alphabet  $A = \{0, 1, \dots, a-1\}$  with length k+1. Note that  $\tau_0(n) = b$  by definition, thus the original unit interval [0, 1) is labeled as b. As the approximation order k increases, each segment is divided into two segments:

(3.3) 
$$[\beta(n)]_k = \left[\beta(n|^k)\right]_k \rightsquigarrow \left[\beta(n|^k)\right]_{k+1} \oplus \left[\beta(n|^k+2^k)\right]_{k+1}.$$

Accordingly, the refinements  $\left[\beta(n^k)\right]_{k+1}$  and  $\left[\beta(n^k+2^k)\right]_{k+1}$  are labeled respectively as

(3.4)  
$$L_{k+1}(n|^{k}) = \tau_{0}(n) \cdots \tau_{k}(n) \left[\frac{\tau_{k}(n)}{2}\right]$$
$$L_{k+1}(n|^{k} + 2^{k}) = \tau_{0}(n) \cdots \tau_{k}(n) \left[\frac{a + \tau_{k}(n)}{2}\right]$$

by Theorem 3.2(2).

Consider a free monoid  $A^* = \bigcup_{k=0}^{\infty} A^k$  over A, with the concatenation  $\oplus$  as a multiplication operation, and the empty string  $\varepsilon$  as a unit, where  $A^k = \{x_1 \oplus x_2 \oplus \cdots \oplus x_k \mid x_i \in A\}$  and  $A^0 = \{\varepsilon\}$ . In view of (3.3) and (3.4), we define a substitution  $\zeta : A^* \to A^*$  as

$$\zeta: x \mapsto \left[\frac{x}{2}\right] \oplus \left[\frac{a+x}{2}\right],$$

which simulates the way to repain the segments that the refinement (3.3) causes:

$$\tau_k(n) \mapsto \tau_{k+1}(n|^k) \oplus \tau_{k+1}(n|^k + 2^k) = \left[\frac{\tau_k(n)}{2}\right] \oplus \left[\frac{a + \tau_k(n)}{2}\right].$$

Consequently, the refinement of the unit interval

 $[0,1) \rightsquigarrow [(0.0)_2)_1 \oplus [(0.1)_2)_1 \rightsquigarrow [(0.00)_2)_2 \oplus [(0.01)_2)_2 \oplus [(0.10)_2)_2 \oplus [(0.11)_2)_2 \rightsquigarrow \cdots$ 

induces an orbit of  $\zeta$  with initial state b:

$$b \mapsto \zeta(b) = \begin{bmatrix} \frac{b}{2} \end{bmatrix} \oplus \begin{bmatrix} \frac{a+b}{2} \end{bmatrix} \mapsto \zeta^2(b) = \begin{bmatrix} \frac{[b/2]}{2} \end{bmatrix} \oplus \begin{bmatrix} \frac{a+[b/2]}{2} \end{bmatrix} \oplus \begin{bmatrix} \frac{a+[(a+b)/2]}{2} \end{bmatrix} \oplus \begin{bmatrix} \frac{a+[(a+b)/2]}{2} \end{bmatrix} \mapsto \cdots$$

Summing up, we have the following.

**Proposition 3.4.** The k-th refinement (2.2) of the unit interval is painted in the pattern  $\zeta^k(b)$ .

3.3. A finite state transducer which represents ax + b. Theorem 3.2 (2) is also expressed in terms of a finite state automaton. Let  $T_{a,b} = (A, I, O, \{b\}, \{0\}, \delta)$  be a deterministic finite state transducer, with states A, an input alphabet  $I = \{0, 1\}$ , an output alphabet  $O = \{0, 1\}$ , an initial state b, a final state 0 and a transition function  $\delta : A \times I \to A \times O$ . According to Theorem 3.2 (2), we define the transition rule as

(3.5) 
$$\delta(x,g) = \begin{cases} \left( \left[ \frac{ag+x}{2} \right], g \right), & \text{if } x \text{ is even} \\ \left( \left[ \frac{ag+x}{2} \right], 1-g \right), & \text{if } x \text{ is odd.} \end{cases}$$

**Corollary 3.5.** Given the binary expressions  $(g_p \cdots g_0)_2$  of a natural number n, define a sequence in  $A \times O$ ,

$$(x_{i+1}, c_{i+1}) = \delta(x_i, g_i), \ i = 0, \dots, q-1,$$

where  $x_0 = b$ ,  $q = p + \operatorname{ord}(a) + 1$  and  $g_i = 0$  if i > p. Then we have

$$f(n) = (c_q c_{q-1} \cdots c_0)_2,$$

and the label of the segment [  $\beta(n)$  )<sub>k</sub> is given as

$$L_k(n) = x_0 x_1 \cdots x_k$$

for each k = 0, 1, ...

*Proof.* Let  $(h_q \cdots h_0)_2$  be the binary expression of f(n). By definition, we see  $x_{i+1} = [(ax_i + g_i)/2]$ . Recall that  $\tau_0(n) = b$  for any  $n \in \mathbb{N}$ . It comes from Theorem 3.2 (2) that

$$\tau_i(n) = x_i$$
 and  $h_i = (g_i + x_i)|^1 = \begin{cases} g_i, & \text{if } x_i \text{ is even,} \\ 1 - g_i, & \text{if } x_i \text{ is odd} \end{cases}$ 

hold, hence  $c_i = h_i$ .

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3.4. Minimality of the dynamical system ([0,1), F). As we have seen, for each approximation order k, the approximant  $F_k$  exchanges the segments  $\left[\beta(\nu_j^{(k)})\right]_k$ 's in the refinement (2.2), and hence induces a permutation  $\pi_k$  over  $\{0, 1, \ldots, 2^k - 1\}$ , satisfying

$$F_k(\left[ \beta(\nu_j^{(k)}) \right)_k) = \left[ \beta(\nu_{\pi_k(j)}^{(k)}) \right)_k,$$

or, equivalently,

$$\nu_{\pi_k(j)}^{(k)} = \left(a\nu_j^{(k)} + b\right)\Big|^k.$$

Note that, as the approximation order k increases, we see

(3.6) 
$$\nu_{2j}^{(k+1)} = \nu_j^{(k)} \text{ and } \nu_{2j+1}^{(k+1)} = \nu_j^{(k)} + 2^k$$

by definition of the refinement (2.2). Combining Theorem 3.2 (3) and (3.6), we have

$$\nu_{\pi_{k+1}(2j)}^{(k+1)} = \nu_{\pi_k(j)}^{(k)} = \nu_{2\pi_k(j)}^{(k+1)} \text{ and } \nu_{\pi_{k+1}(2j+1)}^{(k+1)} = \nu_{\pi_k(j)}^{(k)} + 2^k = \nu_{2\pi_k(j)+1}^{(k+1)}$$

whenever  $\tau_k(\nu_i^{(k)})$  is even, and

$$\nu_{\pi_{k+1}(2j)}^{(k+1)} = \nu_{\pi_k(j)}^{(k)} + 2^k = \nu_{2\pi_k(j)+1}^{(k+1)} \text{ and } \nu_{\pi_{k+1}(2j+1)}^{(k+1)} = \nu_{\pi_k(j)}^{(k)} = \nu_{2\pi_k(j)}^{(k+1)}$$

whenever  $\tau_k(\nu_i^{(k)})$  is odd. As a result, we obtain relations between the permutation  $\pi_k$  and  $\pi_{k+1}$ .

**Lemma 3.6.** If the color 
$$\tau_k(\nu_j^{(k)})$$
 of the segment  $\left[\beta(\nu_j^{(k)})\right]_k$  is even, we have  $\pi_{k+1}(2j) = 2\pi_k(j)$  and  $\pi_{k+1}(2j+1) = 2\pi_k(j) + 1$ 

and if  $\tau_k(\nu_j^{(k)})$  is odd,

$$\pi_{k+1}(2j) = 2\pi_k(j) + 1$$
 and  $\pi_{k+1}(2j+1) = 2\pi_k(j).$ 

Let  $\pi_k = \pi_{k,1}\pi_{k,2}\cdots\pi_{k,m}$  be a cycle decomposition, determined uniquely up to the order of  $\pi_{k,i}$ 's. For each cyclic permutation  $\pi_{k,i} = (p_1p_2\cdots p_l)$ , consider the maximum subset  $Q \subset \{p_1,\ldots,p_l\}$  such that  $\tau_k(\nu_q^{(k)})$ is odd if and only if  $q \in Q$ . For a cycle  $\pi_{k,i} = (p_1p_2\cdots p_l)$ , let  $\pi'_{k+1,i}$  be the permutation defined as Lemma 3.6, that is,

$$\pi'_{k+1,i}(2p_j+\eta) = 2\pi_{k,i}(p_j) + \left(\tau_k(\nu_{p_j}^{(k)}) + \eta\right)\Big|^1 = 2p_{j+1} + \left(\tau_k(\nu_{p_j}^{(k)}) + \eta\right)\Big|^1, \quad \eta \in \{0,1\},$$

where l + 1 is understood as 1.

**Proposition 3.7.**  $\pi'_{k+1,i}$  is a cyclic permutation whenever #Q is odd, and  $\pi'_{k+1,i}$  consists of two cyclic permutations whenever #Q is even.

*Proof.* Put  $P_0 = \{2p_1, 2p_2, ..., 2p_l\}$  and  $P_1 = \{2p_1 + 1, 2p_2 + 1, ..., 2p_l + 1\}$ . Consider an orbit  $x_1 = 2p_1, x_2 = \pi'_{k+1,i}(x_1), x_3 = \pi'_{k+1,i}(x_2), ...,$  and a map  $\phi(x_j) = \eta$  if  $x_j \in P_{\eta}$ . Notice that  $x_{sl+j} \in \{2p_j, 2p_j + 1\}$  for j = 1, ..., l and  $s \in \mathbf{N}$  by definition. Then we see  $\phi(x_{sl+j+1}) = \phi(x_{sl+j})$  whenever  $p_j \notin Q$ , and  $\phi(x_{sl+j+1}) = 1 - \phi(x_{sl+j})$  whenever  $p_j \in Q$ , hence

$$x_{sl+j+1} = \pi'_{k+1,i}(x_{sl+j}) = 2p_{j+1} + \phi(x_{sl+j+1})$$

for each  $j = 1, \ldots, l$  and  $s \in \mathbf{N}$ .

Suppose that #Q is even. It is seen that  $\phi(x_{l+1}) = \phi(x_1) = 0$ , hence

$$x_{l+1} = 2p_1 + \phi(x_{l+1}) = 2p_1 = x_1$$

namely, the cycle  $(x_1x_2, \cdots x_l)$  is a factor of  $\pi'_{k+1,i}$ . Considering another orbit  $y_1 = 2p_1 + 1, y_2 = \pi'_{k+1,i}(y_1), \ldots$ , a similar argument shows the cycle  $(y_1y_2\cdots y_l)$  is also a factor of  $\pi'_{k+1,i}$ . Consequently, we obtain a decomposition  $\pi'_{k+1,i} = (x_1x_2\cdots x_l)(y_1y_2\cdots y_l)$ .

Suppose that #Q is odd, then we see  $\phi(x_{l+1}) = 1 - \phi(x_1) = 1$ , hence

$$x_{l+1} = 2p_1 + \phi(x_{l+1}) = 2p_1 + 1 = y_1,$$

showing that  $\pi'_{k+1,i}$  is a cycle of the length 2*l*:

$$\pi'_{k+1,i} = (x_1 x_2 \cdots x_l y_1 y_2 \cdots y_l).$$

We are to find a condition that the permutation  $\pi_k$  induced by the approximant  $F_k$  becomes a cycle in any approximation order k. For a string  $\boldsymbol{w} = w_1 \oplus w_2 \oplus \cdots \oplus w_l \in A^*$ , we put  $S(\boldsymbol{w}) = \sum_{i=1}^l w_i$ . Note that the parity of  $S(\boldsymbol{w})$  equals that of the number of entries  $w_i$ 's which are odd.

In view of Proposition 3.7, the original unit interval [0, 1) must be painted an odd color at least, namely,  $b = \tau_0([0, 1))$  is odd. Then,  $S(\zeta(b))$  also should be odd. We have assumed that a is odd. Putting a = 2m + 1, we have

$$S(\zeta(b)) = \left[\frac{b}{2}\right] + \left[\frac{a+b}{2}\right] = m+1+2\left[\frac{b}{2}\right]$$

as [(b+1)/2] = [b/2] + 1, hence m should be even;  $a \equiv 1 \pmod{4}$ . Conversely, we have the following.

**Theorem 3.8.** The permutation  $\pi_k$  associated with the approximant  $F_k$  is a cyclic permutation in any approximation order k, hence the dynamical system ([0,1), F) is minimal, if and only if  $a \equiv 1 \pmod{4}$  and b is odd.

Proof. The necessity for b being odd and  $a \equiv 1 \pmod{4}$  is shown above. Suppose a = 4m + 1 and b being odd. It comes from Theorem 3.2 (2) that  $S(\zeta(x)) = 2m + 2[x/2]$  if x is even, and  $S(\zeta(x)) = 2m + 1 + 2[x/2]$  if x is odd. Then,  $S(\zeta(\boldsymbol{w}))$  is odd if and only if  $S(\boldsymbol{w})$  is odd for any  $\boldsymbol{w} \in A^*$ . As a result, we see that  $S(\zeta^k(b))$  is odd for any k, hence  $\pi_k$  is a cyclic permutation by Proposition 3.7, which means for any segment  $[x ]_k$ , the orbit  $[x ]_k$ ,  $F_k([x ]_k)$ ,  $F_k^2([x ]_k)$ , ...,  $F_k^{2^k-1}([x ]_k)$  covers [0, 1):

(3.7) 
$$\prod_{i=0}^{2^{k}-1} F_{k}^{i}\left([x]_{k}\right) = [0,1)$$

Take any  $x \in [0, 1)$ . (3.7) shows that, for any  $y \in [0, 1)$ , there exists  $0 \le t \le 2^k - 1$  with  $y \in F^t([x]_k)$ , while  $F^t(x) \in [F^t(x)]_k = F^t([x]_k)$ . Thus we have  $|y - F^t(x)| < 2^{-k}$ . Consequently, the orbit  $\{F^s(x) \mid s \in \mathbf{N}\}$  is dense in [0, 1) for any  $x \in [0, 1)$ , namely the dynamical system ([0, 1), F) is minimal.

This result is contrast to Keane condition[6][16]: no left endpoint of any segment is mapped to another ones, which brings the minimality of the dynamics of interval exchange maps. In our case, the left endpoint  $x|_k$  of any segment  $[x]_k$  is always mapped to another ones by  $F_k$ .

**Example 3.4.1** (The case 3x+1). This is the original case of Collatz. The substitution  $\zeta$  and the transducer  $T_{3,1}$  are illustrated in FIGURE 1 and 2 respectively. As  $3 \not\equiv 1 \pmod{4}$ , the dynamics ([0,1), F) is not minimal. Indeed, TABLE 1 shows that permutations  $\pi_k$ 's associated with  $F_k$ 's are decomposed in two cycles. A typical orbit of  $F_5$ , the approximation order k = 5, is illustrated in FIGURE 5

k	$\mid \pi_k$
0	(0)
1	(0)(1)
2	$(0\ 2)(1\ 3)$
3	$(0\ 4\ 1\ 5)(2\ 7\ 3\ 6)$
4	$(0 \ 8 \ 2 \ 11 \ 1 \ 9 \ 3 \ 10)(4 \ 14 \ 6 \ 12 \ 5 \ 15 \ 7 \ 13)$
	TABLE 1. Permutations $\pi_k$ associated

TABLE 1. Permutations  $\pi_k$  associate with 3x + 1



FIGURE 1. The color pattern of refinement induced by substitution  $\zeta$ 

**Example 3.4.2** (The case 5x + 1). The substitution  $\zeta$  and the transducer  $T_{5,1}$  are illustrated in FIGURE 3 and 4 respectively. As  $5 \equiv 1 \pmod{4}$ , the conjugacy F of f(x) = 5x + 1 induces a minimal dynamical system on [0, 1). Actually, TABLE 2 shows that permutations  $\pi_k$ 's associated with  $F_k$ 's are always cyclic. A typical orbit of  $F_5$ , the approximation order k = 5, is illustrated in FIGURE 6.



FIGURE 3. The color pattern of refinement induced by substitution  $\zeta$ 



FIGURE 5. Approximant  $F_5$  for 3x + 1 and a typical orbit



FIGURE 2. Transitive diagram of  $T_{3,1}$ 

k	$\pi_k$
0	(0)
1	$(0\ 1)$
2	$(0\ 2\ 1\ 3)$
3	$(0\ 4\ 3\ 7\ 1\ 5\ 2\ 6)$
4	$\left  \begin{array}{cccccccccccccccccccccccccccccccccccc$

TABLE 2. Permutations  $\pi_k$  associated with 5x + 1



FIGURE 4. Transitive diagram of  $T_{5,1}$ 



FIGURE 6. Approximant  $F_5$  for 5x + 1 and a typical orbit

#### 4. Remarks

Notice that F preserves the Lebesgue measure  $\mu$  on [0,1) by Proposition 2.5, therefore, it will be a worthwhile question whether the minimality implies the ergodicity in our case. However, I have no statement about the ergodicity of F at this moment.

In [3] and [4], we have discussed the van der Corput embedding of the Collatz procedure

$$G(x) = \begin{cases} 2x, & x \in [0, 1/2), \\ F(x) = \lim_{k \to \infty} \beta(3\beta^{-1}(x|_k) + 1), & x \in [1/2, 1). \end{cases}$$

Then, to solve the original Collatz conjecture 1.1, we are to prove the statement: 'for each  $n \in \mathbf{N}$ , there exists a finite number t with  $G^t(\beta(n)) = 1/2$ .' Note that G does not preserve intervals, and so the Lebesgue measure. The finite bit approximation

$$G_k(x) = \begin{cases} x + x|_k, & \text{for } x \in [0, 1/2), \\ F_k(x), & \text{for } x \in [1/2, 1) \end{cases}$$

leads us to a mild version of the Collatz conjecture, a finite combinatorial problem proposed in [4], of which research is in progress.

**Problem 4.1**  $(3x + 1 \text{ problem on } G_k)$ . Show that for any  $x \in [0, 1)$ , there exists  $t \in \mathbb{N}$  such that

$$G_k^t(x) \in [0]_k \cup [\beta(1)]_k = [0, 1/2^k] \cup [1/2, 1/2 + 1/2^k].$$

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