

Interval Preserving Map Approximation of $3x + 1$ Problem

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1. Introduction

The well-known and still unsolved $3x + 1$ problem was firstly proposed by Lothar Collatz in 1930's, who had great interest in representation of integer functions by directed graphs.

Conjecture 1.1. *Consider a map $f : \mathbf{N} \rightarrow \mathbf{N}$ such that*

$$f(n) = \begin{cases} 3n + 1, & \text{if } n \text{ is odd,} \\ n/2, & \text{if } n \text{ is even.} \end{cases}$$

Then for each natural number n , there exists a finite number t such that $f^t(n) = \underbrace{f \circ f \circ \cdots \circ f}_{t\text{-times}}(n) = 1$.

In [2], by introducing the *reverse binary embedding of natural numbers* into $[0, 1)$, we obtained a graph \mathfrak{K} of the Collatz function f , which is a Cantor set generated by an iterated function system. Our interest is in the dynamics on the Cantor set \mathfrak{K} presented by the iteration of the Collatz procedure f . In this note, we advance the analysis of the dynamics. We introduce interval preserving maps on $[0, 1)$, approximating the dynamics of the Collatz procedure, and consider a $3x + 1$ problem of ‘finite bit’ version.

2. Reverse binary embedding of natural numbers and a conjugacy of Collatz procedure

Definition 2.1. *Let $n = a_k \cdot 2^k + a_{k-1} \cdot 2^{k-1} + \cdots + a_0$ ($= (a_k a_{k-1} \cdots a_0)_2$ for short) be a binary expansion of a natural number n . The reverse binary embedding β of n is given by*

$$\beta(n) = \frac{a_0}{2} + \frac{a_1}{2^2} + \cdots + \frac{a_k}{2^{k+1}} = (0.a_0 a_1 \cdots a_k)_2.$$

By definition, $\beta : \mathbf{N} \rightarrow [0, 1)$ is one-to-one and \mathbf{N} is densely embedded into $[0, 1)$. Moreover, even and odd numbers are embedded in $[0, 1/2)$ and $[1/2, 1)$ respectively. We note that the binary expression $x = (0.b_1 b_2 \cdots)_2$ of a real number $x \in [0, 1)$ has ambiguity. If there exists $k \in \mathbf{N}$ such that $b_l = 1$ for all $l \geq k$, then $x = (0.b_1 b_2 \cdots b_{k-1})_2 + 2^{-k+1}$, hence x has another binary expression. Thus we always assume that any binary expression $x = (0.b_1 b_2 \cdots)_2$ has infinitely many 0's in the sequence b_j 's, including the case that x has a finite expression $x = (0.b_1 b_2 \cdots b_k)_2$. For a real number $x \in [0, 1)$ and $k \in \mathbf{N}$, $x|_k$ denotes a cut off of x at $(-k - 1)$ -th order in the binary expression;

$$x|_k = (0.b_1 b_2 \cdots b_k)_2 = \sum_{j=1}^k \frac{b_j}{2^j}, \quad \text{if } x = (0.b_1 b_2 \cdots)_2 = \sum_{j=1}^{\infty} \frac{b_j}{2^j}.$$

We use a notation $[x]_k$ as a segment $[x|_k, x|_k + \frac{1}{2^k})$, then we have a natural decomposition of segments

$$(2.1) \quad [(0.b_1 b_2 \cdots b_k)_2]_k = [(0.b_1 b_2 \cdots b_k 0)_2]_{k+1} \oplus [(0.b_1 b_2 \cdots b_k 1)_2]_{k+1},$$

where $[a, b] \oplus [b, c]$ stands for a division of an interval $[a, c]$ at b . For a natural number n , we put $\text{ord}(n) = \lceil \log_2 n \rceil$, where $\lceil x \rceil$ stands for an integer not greater than x , and for a binary expression $n = (a_l a_{l-1} \cdots a_0)_2$, $n|^{k}$ denotes an *upper cut off* of n at k -th order;

$$n|^{k} = (a_{k-1} a_{k-2} \cdots a_0)_2 \equiv n \pmod{2^k}.$$

Apparently we have

Lemma 2.2. *For $x, y \in [0, 1)$ and $m, n, k \in \mathbf{N}$,*

- (1) $x \in [x]_k$.
- (2) $[x]_k \cap [y]_k \neq \emptyset$ holds if and only if $x|_k = y|_k$, and hence $[x]_k = [y]_k$.
- (3) if $x \neq y$, there exists $l \in \mathbf{N}$ such that $[x]_l \cap [y]_l = \emptyset$.
- (4) $[x]_{k+1} \subsetneq [x]_k$ and $\bigcap_{k=1}^{\infty} [x]_k = \{x\}$.
- (5) $\beta(m)|_k = \beta(m|_k)$.
- (6) Thus $[\beta(m)]_k = [\beta(n)]_k$ if and only if $m|_k = n|_k$.

To observe the dynamics of the Collatz procedure, we consider a map $F : [0, 1) \rightarrow [0, 1)$ conjugate to f , that is, $F \circ \beta(n) = \beta \circ f(n)$ holds for any natural number n . The embedding β brings not only a well-defined map $F : \beta(\mathbf{N}) \rightarrow \beta(\mathbf{N})$, but also an extension of F on $[0, 1)$ as follows.

Since $\beta(2n) = \beta(n)/2$, we have $\beta(n) = \beta \circ f(2n) = F \circ \beta(2n) = F(\beta(n)/2)$, hence $F(x) = 2x$ for $x \in [0, 1/2) \cap \beta(\mathbf{N})$. Thus we extend F on $[0, 1/2)$ by $F(x) = 2x$. In order to extend F on $[1/2, 1)$, we need the followings.

Lemma 2.3. *For any natural number m, n , $[\beta(3m+1)]_k = [\beta(3n+1)]_k$ holds if and only if $[\beta(m)]_k = [\beta(n)]_k$.*

Proof. By Lemma 2.2, $[\beta(m)]_k = [\beta(n)]_k$ means $m \equiv n \pmod{2^k}$. Then $3m+1 \equiv 3n+1 \pmod{2^k}$, and hence $[\beta(3m+1)]_k = [\beta(3n+1)]_k$. The converse also holds since 3 and 2^k are prime to each other. \square

Consider a sequence $n_k = \beta^{-1}(x|_k)$, $k = 1, 2, \dots$, then we see $\beta(n_k) = x|_k = (x|_{k+1})|_k = \beta(n_{k+1})|_k$. By Lemma 2.2 (5), we see $n_{k+1} \equiv n_k \pmod{2^k}$, hence $3n_{k+1}+1 \equiv 3n_k+1 \pmod{2^k}$. (2.1) shows

$$[\beta(3n_{k+1}+1)]_{k+1} \subsetneq [\beta(3n_{k+1}+1)]_k = [\beta(3n_k+1)]_k,$$

and then $\bigcap_{k=1}^{\infty} [\beta(3n_k+1)]_k$ consists of a unique point, expressed as $\lim_{k \rightarrow \infty} \beta(3n_k+1)$. Thus we define a conjugacy $F : [0, 1) \rightarrow [0, 1)$ of the Collatz procedure f ;

$$(2.2) \quad F(x) = \begin{cases} 2x, & x \in [0, 1/2), \\ \lim_{k \rightarrow \infty} \beta(3\beta^{-1}(x|_k) + 1), & x \in [1/2, 1). \end{cases}$$

We see that for any odd number n , $\beta(n)|_k = \beta(n) \in [1/2, 1)$ by taking $k > \text{ord}(n)$, therefore $F(\beta(n)) = \beta(3n+1) = \beta(f(n))$. We also see that $[x]_k = [\beta(n_k)]_k$ and $[F(x)]_k = [\beta(3n_k+1)]_k$ as $F(x) \in [\beta(3n_k+1)]_k$ by definition.

Lemma 2.4. $\{\beta(3n+1) \mid n \in \mathbf{N}\}$ is dense in $[0, 1)$.

Proof. As $2 \equiv (-1) \pmod{3}$, we have $2^{2k} \equiv 1 \pmod{3}$ and $2^{2k+1} \equiv (-1) \pmod{3}$. Given $x \in [0, 1)$ and $k \in \mathbf{N}$, consider $m_k = \beta^{-1}(x|_k)$. Suppose $m_k \equiv 0 \pmod{3}$, then there exists $n \in \mathbf{N}$ with $m_k + 2^{2k} = 3n+1$, hence

$$\beta(3n+1) = \beta(m_k + 2^{2k}) = x|_k + \frac{1}{2^{2k+1}} \in [x]_k.$$

Similarly when $m_k \equiv 1 \pmod{3}$ (resp. $m_k \equiv (-1) \pmod{3}$), we find $n \in \mathbf{N}$ with $m_k = 3n+1$ (resp. $m_k + 2^{2k+1} = 3n+1$), showing that $\beta(3n+1) \in [x]_k$. As a result, for any $x \in [0, 1)$ and $k \in \mathbf{N}$, there exists $n \in \mathbf{N}$ with $|x - \beta(3n+1)| < 2^{-k}$. \square

The following proposition implies that F is approximated by interval exchange maps.

Proposition 2.5. *For any odd number n and any natural number k , the conjugacy F is a right continuous bijection*

$$F : [\beta(n)]_k \rightarrow [\beta(3n+1)]_k.$$

Proof. Consider $x, y \in [\beta(n)]_k$ with $x \neq y$, then we find a natural number $l > k$ with $[x]_l \cap [y]_l = \emptyset$. Applying Lemma 2.2 and 2.3 for $x|_l$ and $y|_l$, we see $[\beta(3\beta^{-1}(x|_l) + 1)]_l \cap [\beta(3\beta^{-1}(y|_l) + 1)]_l = \emptyset$, showing $F(x) \neq F(y)$, that is, F is injective on $[\beta(n)]_k$. Lemma 2.3 and 2.4 means that $[\beta(n)]_k \cap \beta(\mathbf{N})$ is embedded in $[\beta(3n + 1)]_k$ densely by F . Then for any $y \in [\beta(3n + 1)]_k$ and natural number $l > k$, we choose $m_l \in \mathbf{N}$ such that $\beta(m_l) \in [\beta(n)]_k$ and $\beta(3m_l + 1) \in [y]_l$. As $\beta(3m_{l+1} + 1) \in [y]_{l+1} \subset [y]_l$, we see $\beta(3m_{l+1} + 1)|_l = \beta(3m_l + 1)|_l$, hence $\beta(m_{l+1})|_l = \beta(m_l)|_l$ by Lemma 2.3. Thus $[\beta(m_{l+1})]_{l+1} \subsetneq [\beta(m_l)]_l$. Let x be a unique accumulation point $\{x\} = \bigcap_{l>k}^{\infty} [\beta(m_l)]_l$, then $x|_l = \beta(m_l)|_l = \beta(m_l|^l)$. We see $y|_l = \beta(3m_l + 1)|_l = \beta(3m_l|^l + 1)|_l$ as $3m_l + 1 \equiv 3m_l|^l + 1 \pmod{2^l}$, hence

$$\beta(3\beta^{-1}(x|_l) + 1) = \beta(3m_l|^l + 1) \in [y]_l.$$

Taking $l \rightarrow \infty$, we have $F(x) = y$, showing that F is a surjection $[\beta(n)]_k \rightarrow [\beta(3n + 1)]_k$.

For any $x \in [\beta(n)]_k$ and $l > k$, we see $F(x) \in [\beta(3\beta^{-1}(x|_l) + 1)]_l$ by the definition of F . We also have $F(w) \in [\beta(3\beta^{-1}(w|_l) + 1)]_l = [\beta(3\beta^{-1}(x|_l) + 1)]_l$ whenever $w \in [x]_l$, that is, $|F(w) - F(x)| < 2^{-l}$ if $w \in [x]_l$. Therefore F is right continuous. \square

Note that for any natural number n , $[\beta(n)]_k = [\beta(n), \beta(n) + 2^{-k})$ if $k > \text{ord}(n)$, thus $w \in [\beta(n)]_k$ means $w \geq \beta(n)$. Therefore we state the right continuity only in Proposition 2.5. Indeed, for example, we see

$$\lim_{\substack{w \rightarrow (0.11)_2 \\ w < (0.11)_2}} F(w) = (0.001)_2 \neq (0.0101)_2 = F((0.11)_2).$$

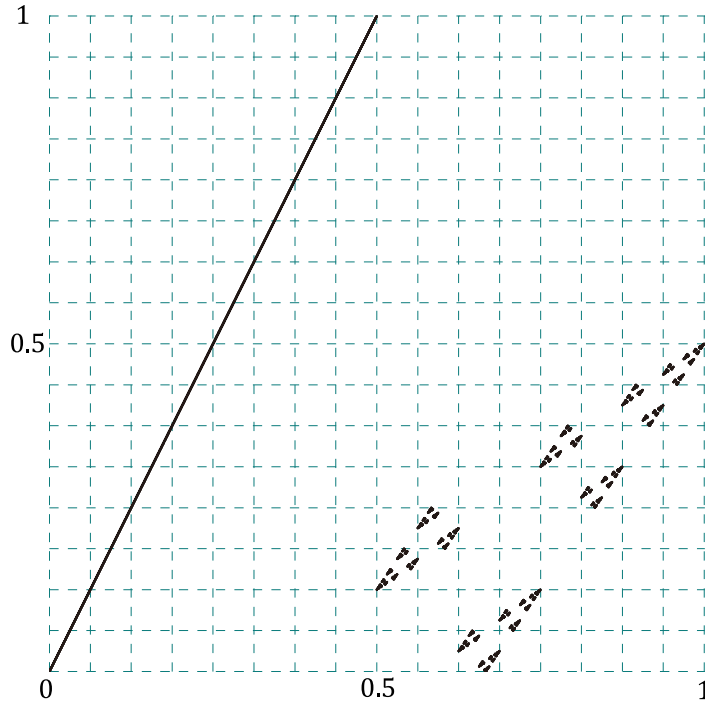


FIGURE 1. The graph of F conjugate to the Collatz procedure

3. Piecewise linear approximant of the conjugacy F

In view of Proposition 2.5, we construct a sequence of piecewise linear maps F_k approximating the conjugacy F : for each $k \in \mathbf{N}$, we define the k -th approximant F_k as

$$\begin{aligned} F_k(x) &= \begin{cases} 2x, & \text{for } x \in [0, 1/2), \\ x - x|_k + F(x|_k)|_k, & \text{for } x \in [1/2, 1), \end{cases} \\ &= \begin{cases} 2x, & \text{for } x \in [0, 1/2), \\ x - x|_k + \beta(3\beta^{-1}(x|_k) + 1)|_k, & \text{for } x \in [1/2, 1). \end{cases} \end{aligned}$$

By definition, F_k is right continuous. Since $w|_k = \beta(n)|_k$ if $w \in [\beta(n)|_k)$ and $\beta(3n|_k + 1)|_k = \beta(3n + 1)|_k$, we see

$$F(w) = w - \beta(n)|_k + \beta(3n + 1)|_k, \quad \text{for } w \in [\beta(n)|_k).$$

Thus $F_k([\beta(n)|_k)) = [\beta(3n + 1)|_k)$ for any odd number n , compatible with Proposition 2.5. As a result, we see $F(x), F_k(x) \in [\beta(3n + 1)|_k)$ for any $x \in [\beta(n)|_k)$, showing that

$$|F(x) - F_k(x)| < 2^{-k}$$

holds for any $x \in [1/2, 1)$. Therefore the sequence F_k , $k = 1, 2, \dots$ approximates F uniformly on $[1/2, 1)$.

In the following subsections, we investigate the approximants F_k 's to extract dynamical characteristics of the map $3x + 1$.

3.1. Behavior of carries from lower to upper bits. For $n, k \in \mathbf{N}$, we define an integer valued function

$$\tau_k(n) = \left\lfloor \frac{3n|_k + 1}{2^k} \right\rfloor.$$

The function τ_k describes the number of bits carried in the calculation $3n + 1$ at k -th bit, as follows.

Proposition 3.1. *Given an odd number n and take $k \in \mathbf{N}$, then we have*

- (1) $\tau_k(n) \in \{0, 1, 2\}$.
- (2) For binary expressions $n = (a_l \cdots a_0)_2$ and $3n + 1 = (c_l \cdots c_0)_2$ with $l \geq \text{ord}(n) + 2$, we have

$$c_k = (a_k + \tau_k(n))|_1 \quad \text{for } k \leq l.$$

- (3)
$$\tau_{k+1}(n|_k) = \begin{cases} 0, & \text{if } \tau_k(n) = 0, 1, \\ 1, & \text{if } \tau_k(n) = 2, \end{cases} \quad \tau_{k+1}(n|_k + 2^k) = \begin{cases} 1, & \text{if } \tau_k(n) = 0, \\ 2, & \text{if } \tau_k(n) = 1, 2. \end{cases}$$

- (4) If $\tau_k(n) = 0$ or 2 ,

$$\left[F(\beta(n|_k)) \right]_k = \left[F(\beta(n|_k)) \right]_{k+1} \oplus \left[F(\beta(n|_k + 2^k)) \right]_{k+1}$$

holds, and if $\tau_k(n) = 1$,

$$\left[F(\beta(n|_k)) \right]_k = \left[F(\beta(n|_k + 2^k)) \right]_{k+1} \oplus \left[F(\beta(n|_k)) \right]_{k+1}$$

holds.

Proof. (1) As $n|_k \leq 2^k - 1$, we have

$$0 \leq \left\lfloor \frac{3n|_k + 1}{2^k} \right\rfloor \leq \left\lfloor \frac{3 \cdot 2^k - 2}{2^k} \right\rfloor = \left\lfloor 3 - \frac{1}{2^{k-1}} \right\rfloor \leq 2.$$

(2) By definition, we see $n|^{k+1} = (a_k a_{k-1} \cdots a_0)_2 = a_k \cdot 2^k + n|^{k+1}$ and $3n|^{k+1} + 1 = \tau_k(n) \cdot 2^k + (3n + 1)|^{k+1}$. We also see $(3n + 1)|^{k+1} = c_k \cdot 2^k + (3n + 1)|^{k+1}$ while

$$(3.1) \quad (3n + 1)|^{k+1} = (3n|^{k+1} + 1)|^{k+1} = \left(3(a_k \cdot 2^k + n|^{k+1}) + 1 \right) \Big|^{k+1} \\ = \left((3a_k + \tau_k(n)) \cdot 2^k + (3n + 1)|^{k+1} \right) \Big|^{k+1} = (3a_k + \tau_k(n)) \Big|^{k+1} \cdot 2^k + (3n + 1)|^{k+1}.$$

Therefore $c_k = (3a_k + \tau_k(n)) \Big|^{k+1} = (a_k + \tau_k(n)) \Big|^{k+1}$.

(3) The case $\tau_{k+1}(n)$ is shown by definition. We also see $3n|^{k+1} + 1 = \tau_k(n) \cdot 2^k + r$ where $r = (3n|^{k+1} + 1) < 2^k$. Then we have

$$2^{k+1} < 3(n|^{k+1} + 2^k) + 1 = (\tau_k(n) + 3) \cdot 2^k + r < (\tau_k(n) + 4) \cdot 2^k.$$

When $\tau_k(n) = 0$, it holds that $2^{k+1} < 3(n|^{k+1} + 2^k) + 1 < 2^{k+2}$, hence $\tau_{k+1}(n|^{k+1} + 2^k) = 1$. When $\tau_k(n) = 1, 2$, we have

$$2^{k+2} = (1 + 3) \cdot 2^k \leq (\tau_k(n) + 3) \cdot 2^k + r < (2 + 4) \cdot 2^k = 3 \cdot 2^{k+1},$$

hence $\tau_{k+1}(n|^{k+1} + 2^k) = 2$.

(4) Taking n for $n|^{k+1} + a_k \cdot 2^k$, it comes from (3.1) that

$$\beta(3(n|^{k+1} + a_k \cdot 2^k) + 1) \Big|_{k+1} = \beta(3n|^{k+1} + 1) \Big|_k + \frac{1}{2^{k+1}} \beta(3a_k + \tau_k(n)) \Big|_1.$$

When $\tau_k(n) = 0$ or 2 , we have $(3a_k + \tau_k(n)) \Big|^{k+1} = a_k$, and when $\tau_k(n) = 1$, we have $(3a_k + \tau_k(n)) \Big|^{k+1} = 1 - a_k$, hence the assertion. \square

Since $[\beta(n|^{k+1})]_{k+1} = [\beta(n|^{k+1})]_{k+1} \oplus [\beta(n|^{k+1} + 2^k)]_{k+1}$, Proposition 3.1 (4) means that F flips the images of $[\beta(n|^{k+1})]_{k+1}$ and $[\beta(n|^{k+1} + 2^k)]_{k+1}$ if and only if $\tau_k(n) = 1$.

3.2. Substitution dynamics and automaton associated with the approximant F_k 's. According to Proposition 3.1, the behavior of τ_k 's brings a substitution dynamics and an automaton.

For an odd number n and $k \in \mathbf{N}$, we call n *stationary* at k -th order when $\tau_k(n) = 0$, *exchanging* when $\tau_k(n) = 1$ and *unfixed* when $\tau_k(n) = 2$ respectively. We label each segment $[\beta(n)]_k$ as S , E and U when n is stationary, exchanging and unfixed at k -th order respectively. The original segment $[\beta(1)]_1 = [1/2, 1)$ is labeled U .

Then combining Proposition 3.1 (3) and (4), to increment the approximation order k by 1 causes a division of each segment, and induces a substitution

$$\sigma : S \rightarrow SE \quad E \rightarrow SU \quad U \rightarrow EU,$$

which are mapped by F as

$$F(\sigma) : F(S) \rightarrow F(S)F(E) \quad F(E) \rightarrow F(U)F(S) \quad F(U) \rightarrow F(E)F(U).$$

Accordingly we associate each $\beta(n)$ with a string $\mathbf{w} = w_1 w_2 \cdots$ consists of S, E and U ; we take w_k for S, E or U when $[\beta(n)]_k$ is labeled as S, E or U , in another words, when $\tau_k(n) = 0, 1$, or 2 respectively. Note that $w_k = S$ if $k > \text{ord}(n) + 2$ as $3n + 1 < 2^{\text{ord}(n)+3}$. The sequence \mathbf{w} is given by calculating $\tau_k(n)$ for each k , however, Proposition 3.1 (3) brings us another way to obtain \mathbf{w} directly, without help of τ_k . Actually, suppose that $(0.b_1 b_2 \cdots b_k)_2$ is stationary S at k -th order for instance, then $(0.b_1 b_2 \cdots b_k b_{k+1})_2$ is again stationary S if $b_{k+1} = 0$ and exchanging E if $b_{k+1} = 1$.

Thus we define a deterministic finite state automaton $M = (\{S, E, U\}, \{0, 1\}, \delta, E, \{S\})$ with states S, E, U , an alphabet $\{0, 1\}$ of input symbols, an initial state E , an accepting state S and a transition

function δ given as

$$\begin{aligned}\delta(S, 0) &= S, & \delta(E, 0) &= S, & \delta(U, 0) &= E, \\ \delta(S, 1) &= E, & \delta(E, 1) &= U, & \delta(U, 1) &= U.\end{aligned}$$

The automaton M encodes a binary $(0.b_1b_2\cdots b_k)_2$ to a string $\mathbf{w} = w_1w_2\cdots w_k$ over $\{S, E, U\}$, by scanning each term of $b_1b_2\cdots b_k$ from left to right. Apparently we see the automaton M comes to the final state S whenever M has read binaries of the form $(0.b_1b_2\cdots b_k00)_2$ completely; any input data of the form $b_1\cdots b_k00$ is accepted by M

Conversely, we restore the binary $F(\beta(n)) = \beta(3n + 1)$ from $\beta(n)$ and its encode \mathbf{w} .

Corollary 3.2. *Consider a binary expression $\beta(n) = (0.b_1b_2\cdots b_k)_2$ with $b_{k-1} = b_k = 0$ for any odd number n , and its encode $\mathbf{w} = w_1w_2\cdots w_k$ that the automaton M outputs. Let $c_1c_2\cdots c_k$ be a string consist of 0 and 1, defined as*

$$c_i = \begin{cases} b_i, & \text{if } w_{i-1} \neq E, \\ 1 - b_i, & \text{if } w_{i-1} = E, \end{cases}$$

where we take w_0 for E . Then we have

$$\beta(3n + 1) = (0.c_1c_2\cdots c_k)_2.$$

Proof. Applying Proposition 3.1 (2) for $n = (b_kb_{k-1}\cdots b_1)_2$ and $3n + 1 = (c_kc_{k-1}\cdots c_1)_2$, we have

$$c_i = (b_i + \tau_{i-1}(n)) \bmod 2 = \begin{cases} b_i, & \text{if } \tau_{i-1}(n) = 0, 2, \\ 1 - b_i, & \text{if } \tau_{i-1}(n) = 1, \end{cases}$$

hence the assertion. □

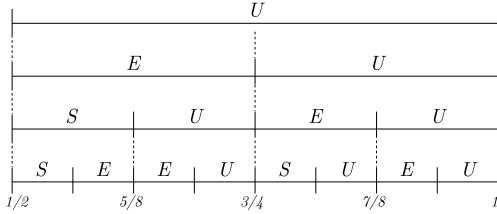


FIGURE 2. Substitution σ on segments

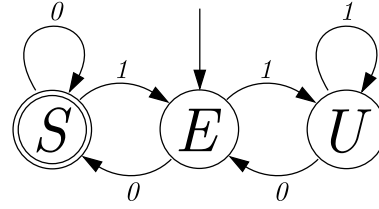


FIGURE 3. Transitive diagram of M

3.3. An attracting property and a $3x + 1$ problem on F_k 's. It easily comes from the definition of F_k that any finitely long binary sequence $\beta(n)$ arrives at a certain k -bit binary sequence by a sufficiently large number of iterations of F_k .

Proposition 3.3. *Consider the k -th approximant F_k . Then for any natural number n , there exists $t \in \mathbf{N}$ that*

$$\text{ord}(\beta^{-1}(F_k^t(\beta(n)))) \leq k.$$

Proof. Denote $n_l = \beta^{-1}(F_k^l(\beta(n)))$. When n_l is odd, decompose $n_l = A_l + B_l$ with $B_l = n_l \bmod 2$ and $A_l = n_l - B_l$. By definition, we see $\beta(n_l)|_k = \beta(B_l)$ and $\beta(n_l) = \beta(A_l) + \beta(B_l)$. Thus $F_k(\beta(n_l)) = \beta(A_l) + F(\beta(B_l))|_k$ by the definition of F_k . As $F_k(\beta(n_l)) \in [0, 1/2)$, we see

$$F_k^2(\beta(n_l)) = 2(\beta(A_l) + F(\beta(B_l))|_k) = \beta(A_l/2) + 2F(\beta(B_l))|_k.$$

As $\text{ord}(A_l/2) = \text{ord}(A_l) - 1$, there exists $t \in \mathbf{N}$ such that $A_t = 0$, hence the assertion. □

Thus the orbit of any finite binary sequence $\beta(n)$ eventually joins that of a k bit one, meaning that we concentrate our attention on k bit binaries to answer the following ‘Collatz-like’ problem on F_k :

Problem 3.4 ($3x + 1$ problem on F_k). *Show that for any natural number n , there exists $t \in \mathbf{N}$ such that*

$$F_k^t(\beta(n)) = 0 \text{ or } 1/2.$$

Note that finite bit approximation causes the fixed point $n = 0$, e.g., for F_3 we have

$$(0.101)_2 \mapsto (0.000)_2,$$

by eliminating binaries smaller than 2^{-3} . Problem 3.4 moderates the original $3x + 1$ problem, where only finite numbers smaller than 2^k are considered. Thus any orbit of $\beta(n)$ are eventually periodic.

4. Interval preserving approximant of the conjugacy F

The approximant F_k behaves as an interval exchange transformation on $[1/2, 1)$, while F_k expands the interval $[0, 1/2)$, which seems to cause some difficulties to close the true nature of the $3x + 1$ problem on F_k . Accordingly, we consider another series of approximants G_k ’s defined as

$$\begin{aligned} G_k(x) &= x - x|_k + F(x|_k)|_k, & \text{for } x \in [0, 1), \\ &= \begin{cases} x + x|_k, & \text{for } x \in [0, 1/2), \\ x - x|_k + \beta(3\beta^{-1}(x|_k) + 1)|_k, & \text{for } x \in [1/2, 1), \end{cases} \end{aligned}$$

which are interval preserving and right continuous, approximating F uniformly on $[0, 1)$;

$$|F(x) - G_k(x)| < 2^{-k} \text{ for all } x \in [0, 1).$$

Since G_k translates only each segment $[\beta(n)|_k)$ to another by definition, the orbit of any point $x \in [0, 1)$ are eventually periodic, described completely by the orbit of the k bit sequence $x|_k$. Thus the corresponding $3x + 1$ problem is stated as follows, which reduces Problem 3.4 to a finite combinatorial one:

Problem 4.1 ($3x + 1$ problem on G_k). *Show that for any $x \in [0, 1)$, there exists $t \in \mathbf{N}$ such that*

$$G_k^t(x) \in [0|_k) \cup [\beta(1)|_k) = \left[0, 1/2^k\right) \cup \left[1/2, 1/2 + 1/2^k\right).$$

Again we note that any $x \in [0|_k)$ is fixed under G_k . In spite of restricting our argument to finite bits, Problem 4.1 seems to be not trivial. If we consider $5x + 1$ map instead of $3x + 1$ for instance, we find lots of periodic orbits as increasing the approximation order k . Problem 4.1 appears to be presenting the unique characteristics of $3x + 1$ map.

To attack the original $3x + 1$ problem, we have to face the fact that for any initial point $\beta(n)$, it seems to exist sufficiently large approximation order k where the orbit of $\beta(n)$ under G_k pass through only stationary segments. Consider $n = 15 = (1111)_2$. Table 1 shows the corresponding orbits under G_7 and G_8 (omitting even numbers). We see that the orbit pass through stationary segments for $k \geq 8$; each segment label is finished at S . Thus the orbit is determined independently of the approximation order $k \geq 8$, meaning that the orbit under G_k coincides with the original one under F . See the graphs of G_1 to G_8 and their orbits of $(0.1111)_2$ in Figures 4 and 5.

Orbit under G_7 :			Orbit under G_8 :		
n	$\beta(n)$	segment label	n	$\beta(n)$	segment label
15	.1111000	<i>UUUU</i> ESS	15	.1111000	<i>UUUU</i> ESSS
23	.1110100	<i>UUUE</i> UES	23	.1110100	<i>UUUE</i> ESSS
35	.1100010	<i>UUES</i> SES	35	.1100010	<i>UUES</i> ESSS
53	.1010110	<i>UEUE</i> UUE	53	.1010110	<i>UEUE</i> UUES
1	.1000000	<i>UESS</i> SSS	5	.1010000	<i>UEUE</i> SSSS
			1	.1000000	<i>UESS</i> SSSS

TABLE 1. Orbits of $n = 15$ under G_7 and G_8

5. A conjecture arisen from a graph symmetric to F

In view of the definition (2.2) of the conjugacy F , we also consider a map H symmetrical about $(1/2, 1/2)$ with F on $\beta(\mathbf{N})$:

$$H(x) = \begin{cases} \lim_{k \rightarrow \infty} \beta(3\beta^{-1}(x|_k) + 1), & x \in [0, 1/2), \\ 2x - 1, & x \in [1/2, 1). \end{cases}$$

In fact, the topological closure of

$$\{(\beta(n), \beta(3n + 1)) \mid n \in \mathbf{N}\}$$

is symmetrical about $(1/2, 1/2)$. For any natural numbers n and k , consider the *one's-complement* $\overline{n|_k}$ of k -bit; $\overline{n|_k} = 2^k - 1 - n|_k$. Then we have $\beta(n|_k) + \beta(\overline{n|_k}) = 1 - 2^{-k}$. It is easily seen that $\overline{(3n + 1)|_k} = \overline{(3\overline{n|_k} + 1)|_k}$. Then we see $1 - 2^{-k} = \beta((3n + 1)|_k) + \beta(\overline{(3n + 1)|_k}) = \beta(3n + 1)|_k + \beta(3\overline{n|_k} + 1)|_k$. As a result, we show

$$\lim_{k \rightarrow \infty} (\beta(\overline{n|_k}), \beta(3\overline{n|_k} + 1)|_k)$$

is symmetrical to $(\beta(n), \beta(3n + 1))$ about $(1/2, 1/2)$.

We also note that H is the conjugacy of the following arithmetic procedure:

$$h(n) = \begin{cases} 3n + 1, & \text{if } n \text{ is even,} \\ (n - 1)/2, & \text{if } n \text{ is odd.} \end{cases}$$

It is easy to find that h has at least three periodic orbits, which pass through $n = 0$, $n = 4$ and $n = 16$ respectively. However observation of the dynamics of H leads us to another conjecture:

Conjecture 5.1. *For each natural number n , there exists a finite number t such that $h^t(n) = 0, 4$ or 16 .*

Note that H is also right continuous on $[0, 1)$, where we are to consider, so to speak, the ‘left continuous version’ of F .

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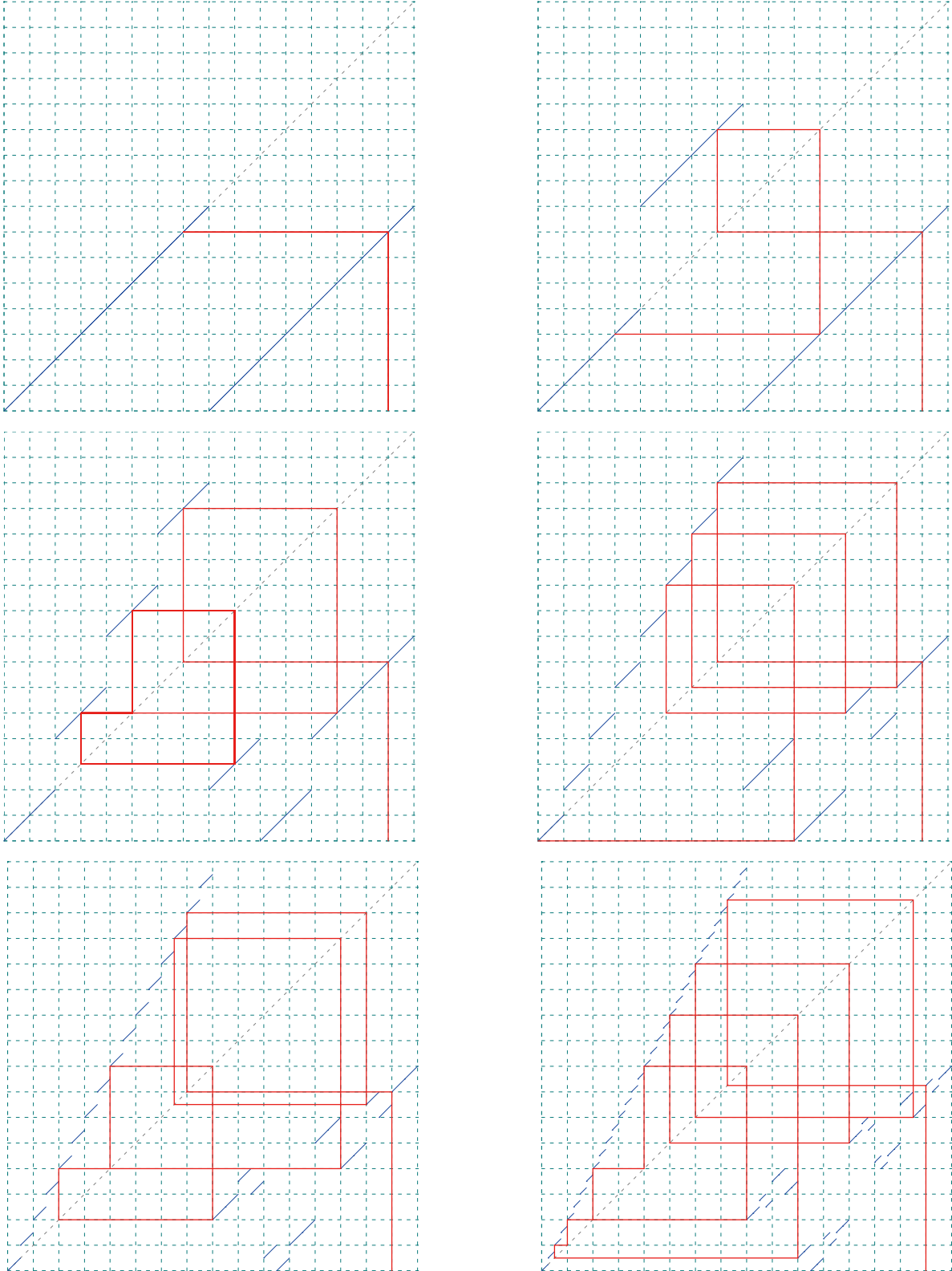


FIGURE 4. Graphs of G_1 to G_6 with their orbit start from $(0.1111)_2$

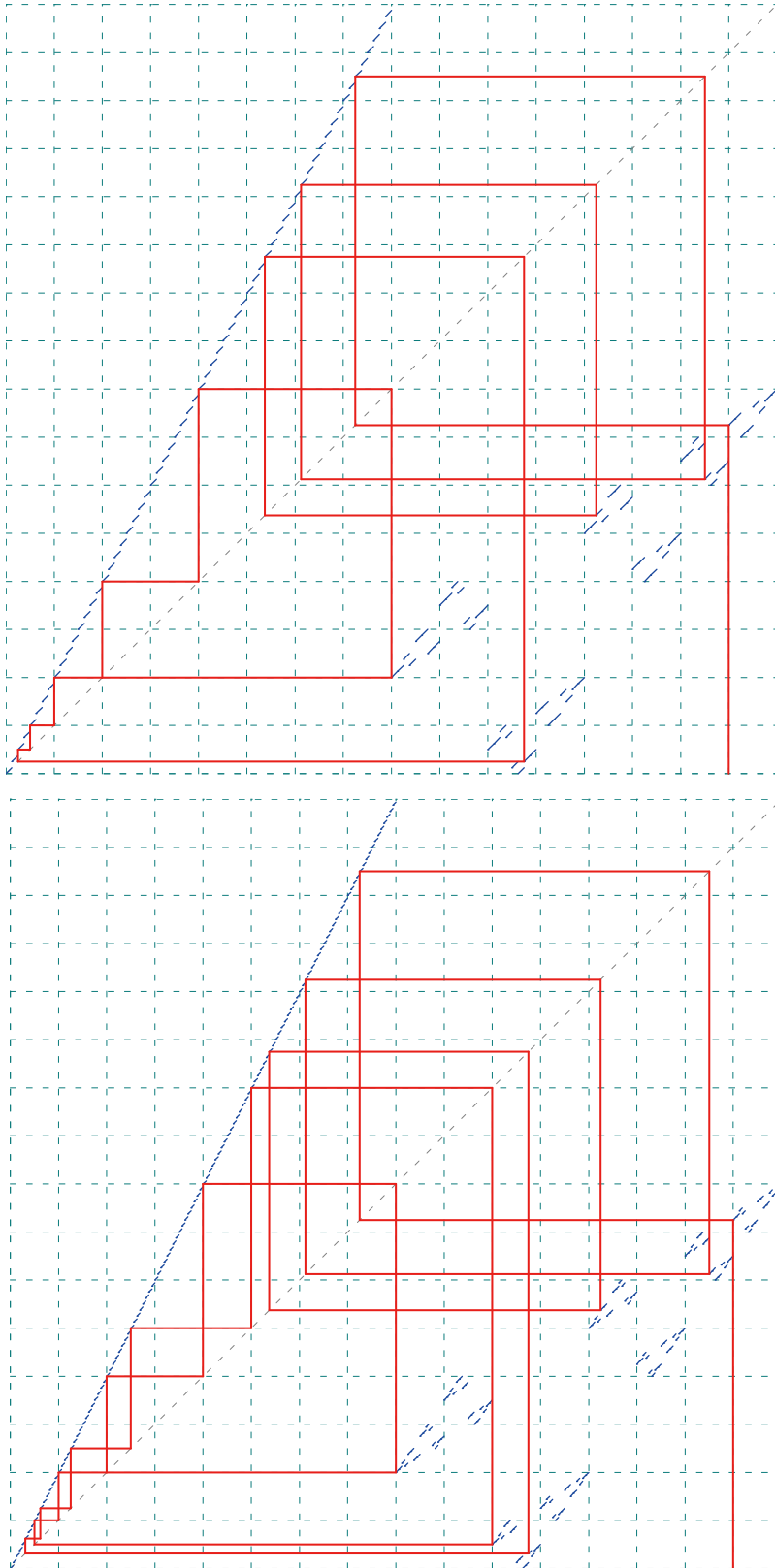


FIGURE 5. Graphs of G_7 and G_8 with their orbit start from $(0.1111)_2$

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